


CSCI 2100

Heap Recap
Binary Trees
Search Trees



Recap

- HW due today
- Next HW - posted after class, due in 1 week
(pen + paper)
- Lab tomorrow

Last "time":

- Trees: binary tree
(height + depth)

- Priority Queues: ADT supporting:

- insert
- get Max
- remove Max

• Heaps:

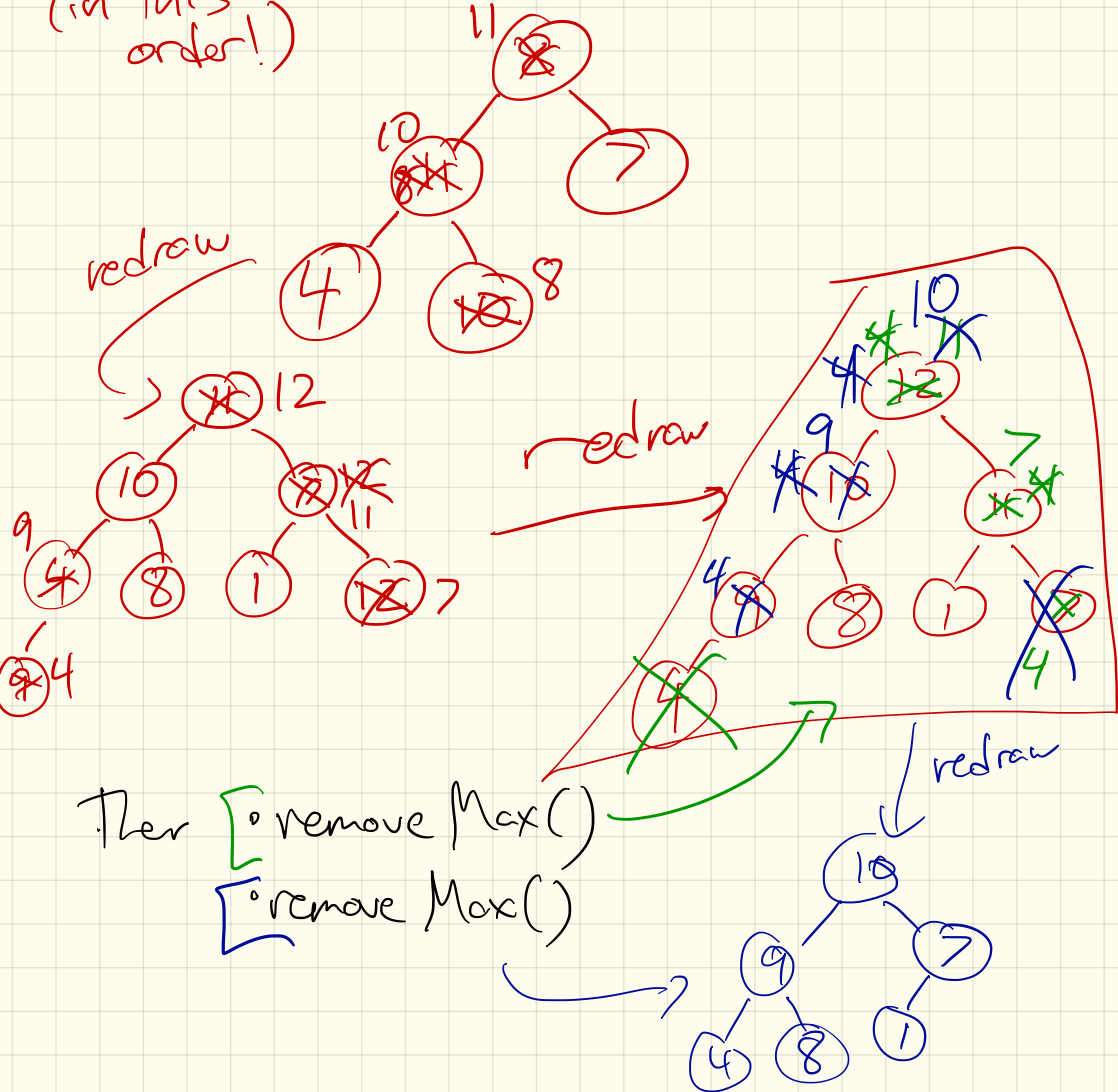
binary tree implementation
of PQ

value at node is
 \geq children's value

all operations are $O(\log n)$

Heap example:

Insert: 8, 11, 3, 4, 10, 1, 12, 9
(in this order!)

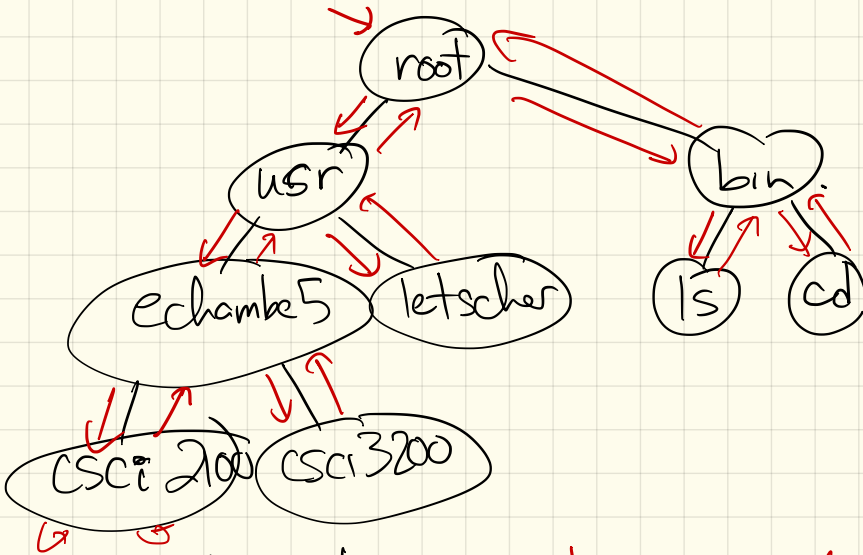


Aside: Tree traversals

3 ways to traverse a tree:
Starting at root:

Preorder (v) =

print \rightarrow visit \checkmark
Preorder (v \rightarrow left)
Preorder (v \rightarrow right)



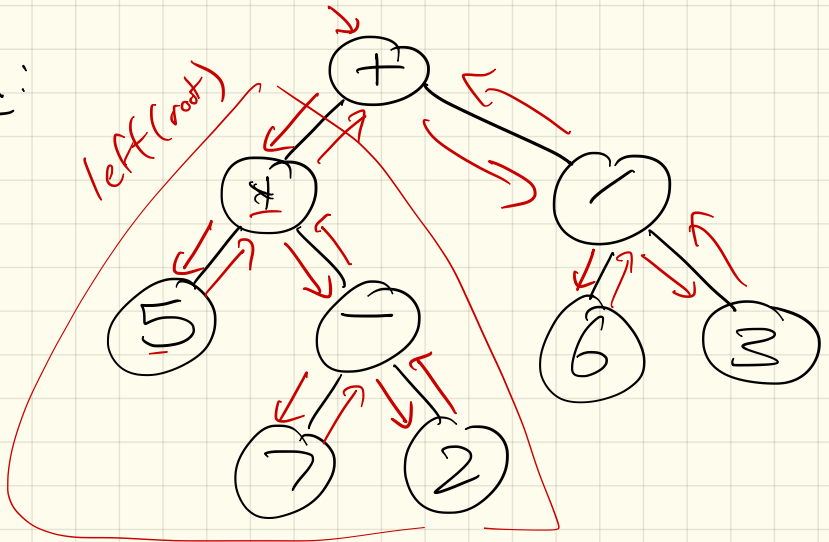
traversal order: root, usr, echambe5,
csci2100, csci3200,
letscher, bin, ls, cd,

(start at root)
In order (v) =

Inorder (v → left)

visit v
Inorder (v → right)

Ex:



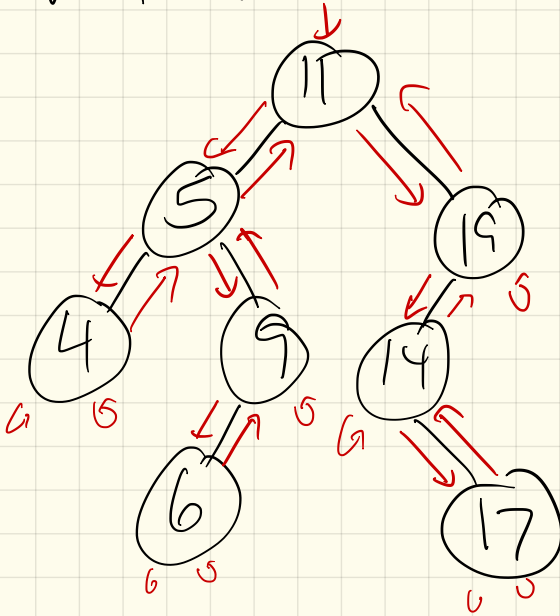
traversal: 5, *, 7, -, 2, +, 6, 3

↑
"middle"

Post order (v) =

post order (v \rightarrow left)
post order (v \rightarrow right)
visit v

Ex:

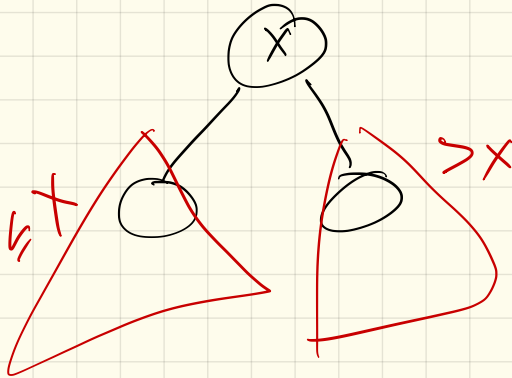


traversal: 4, 6, 9, 5, 17, 14, 19,
11

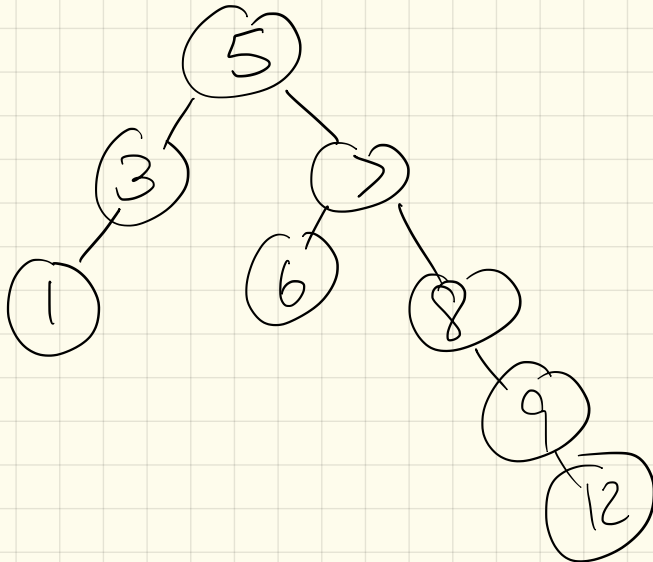
Next: Binary Search Trees

A binary tree where we maintain the following:

- The value at any node is \geq its left child and $<$ its right child.



Ex:

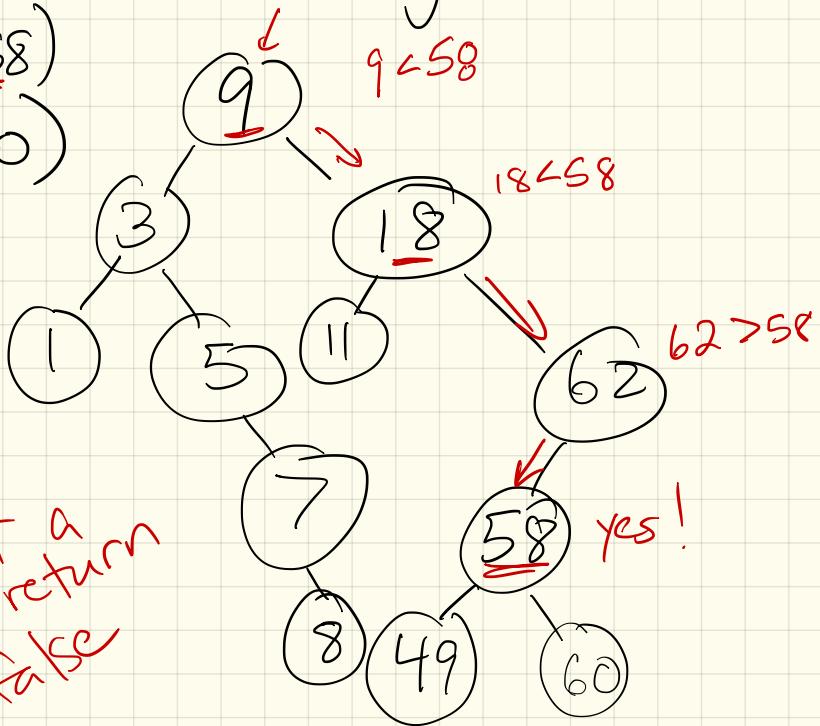


Finding in a BST:

- Check if root = target value
return true
- If root > target
if left != Null
recurse on left child
else return false
- If root < target
recurse on right child

Ex:

find (58)
find (50)

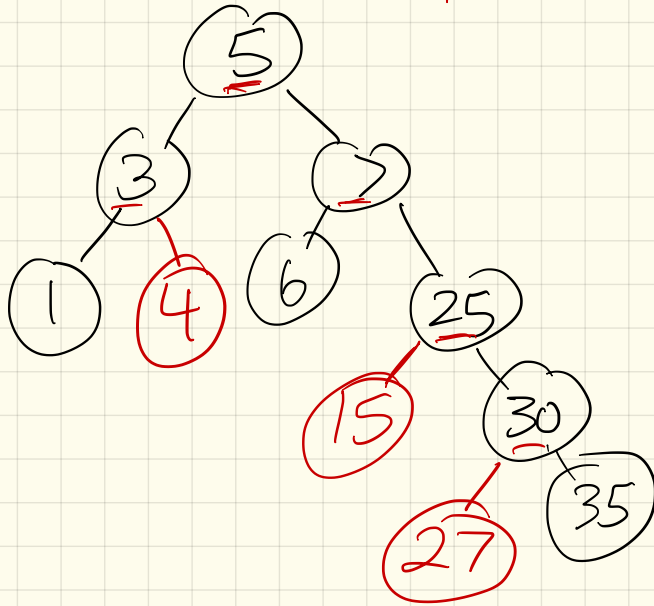


Inserting:

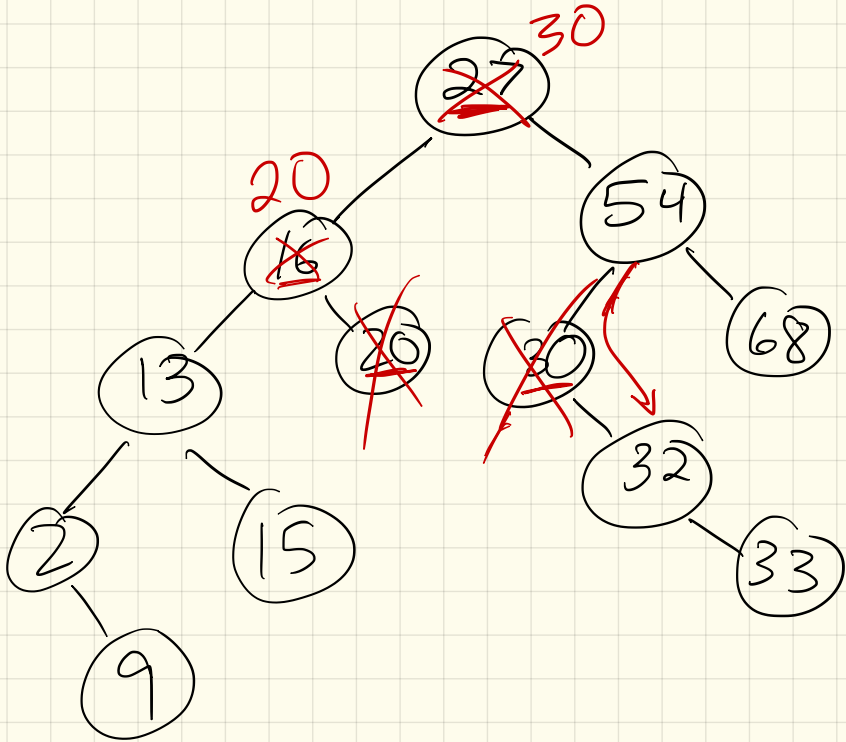
Given a BST, insert is done by finding the (unique) leaf location where the value fits.

Ex: insert (4)
insert (27)
insert (15)

find, but when hit a null, insert the new node



Delete: More Complex



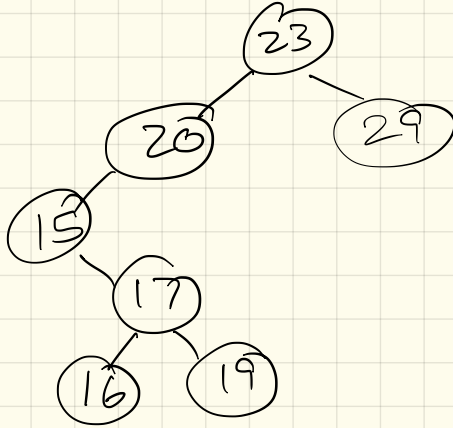
remove(16) : delete 20's node
 copy 20 into 16's node

remove(27) : copy 30 up to root

Cases:

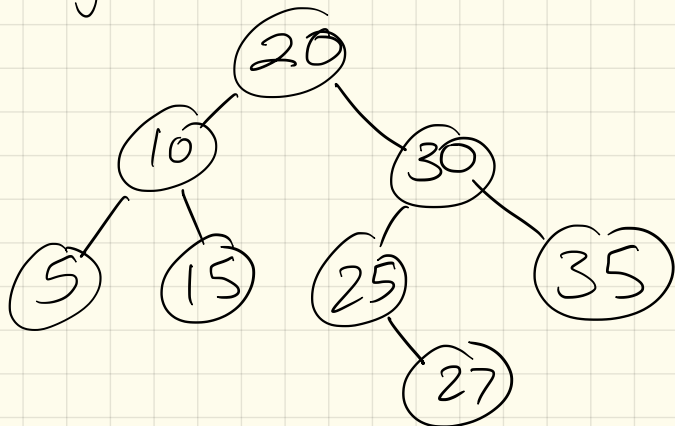
- Target could be a leaf

- Target could have only 1 child: remove(15)



Cases (cont):

• Target could have 2 children



remove (20) :

Remove(x):

Find (x)

If x is a leaf
delete x's node

else if x has one child

can remove x's node
and "promote" its
only child to its spot

Else:

Find smallest value $> x$

(Note: this is next value
in an inorder traversal!)

Copy that value to x's spot
& remove that node,
promoting its right child
in its spot (if it has one)

Note: it will have no
left child

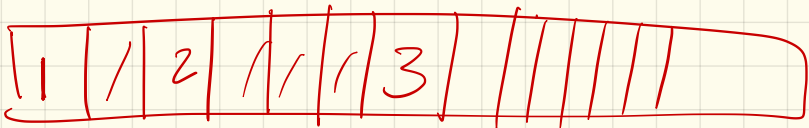
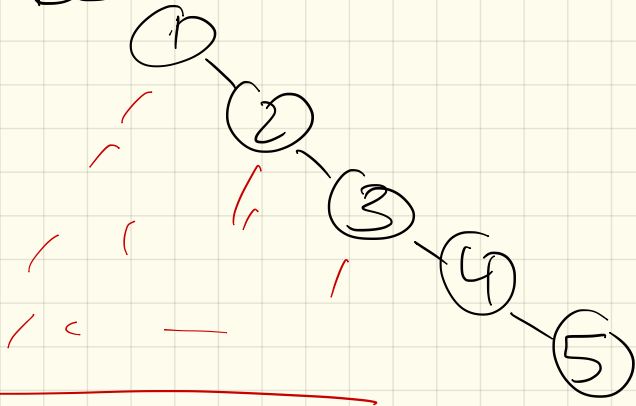
Code:

- Pointer based.

Reason: we need to move around entire subtrees.

Also, tree is not complete!

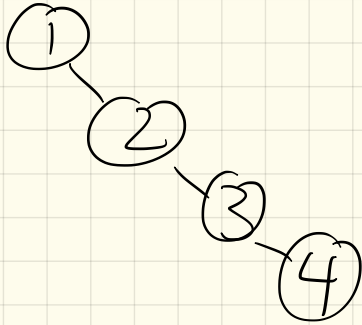
Valid BST:



not space efficient

Note: BSTs are not unique!

Consider



Can you make another BST
with these values?