

CS2100 - AVL Trees

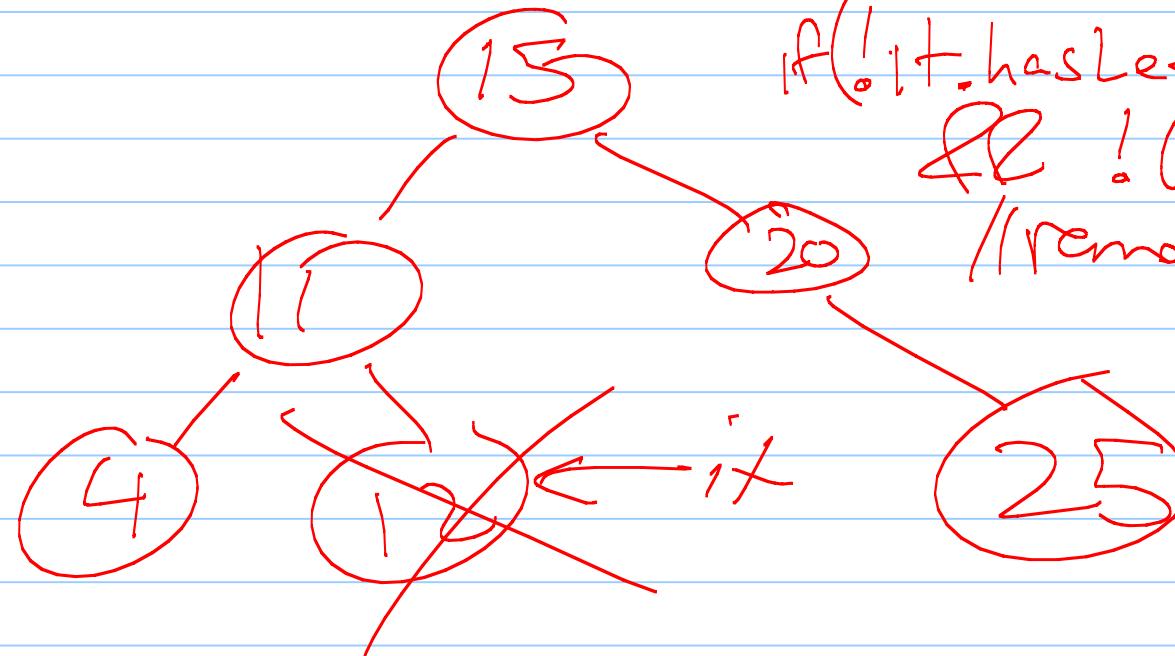
Note Title

10/19/2012

Announcements

~ Scholarship deadline next week

Remove: Cases



remove (12)

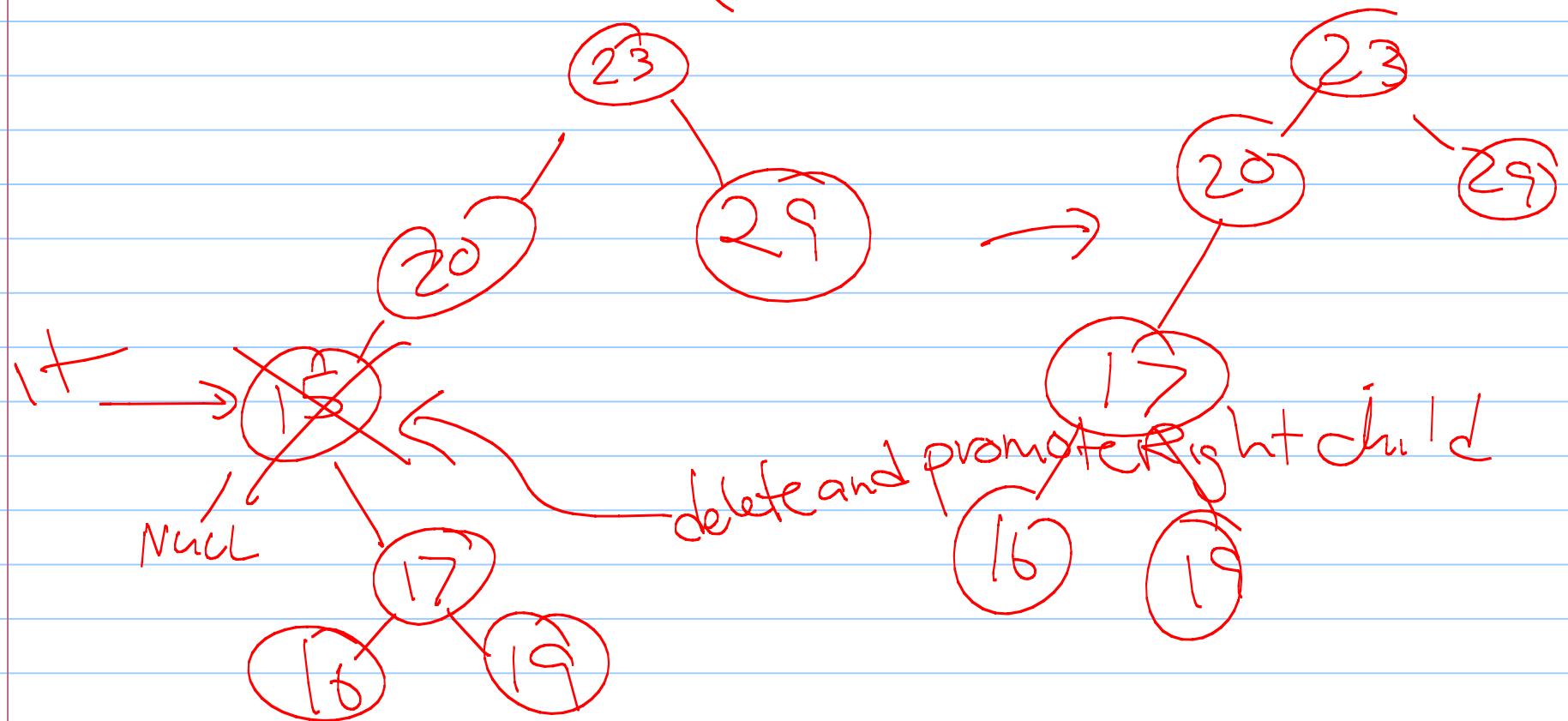
If (!it.hasLeftChild())

OR ! (it.haveRightChild())

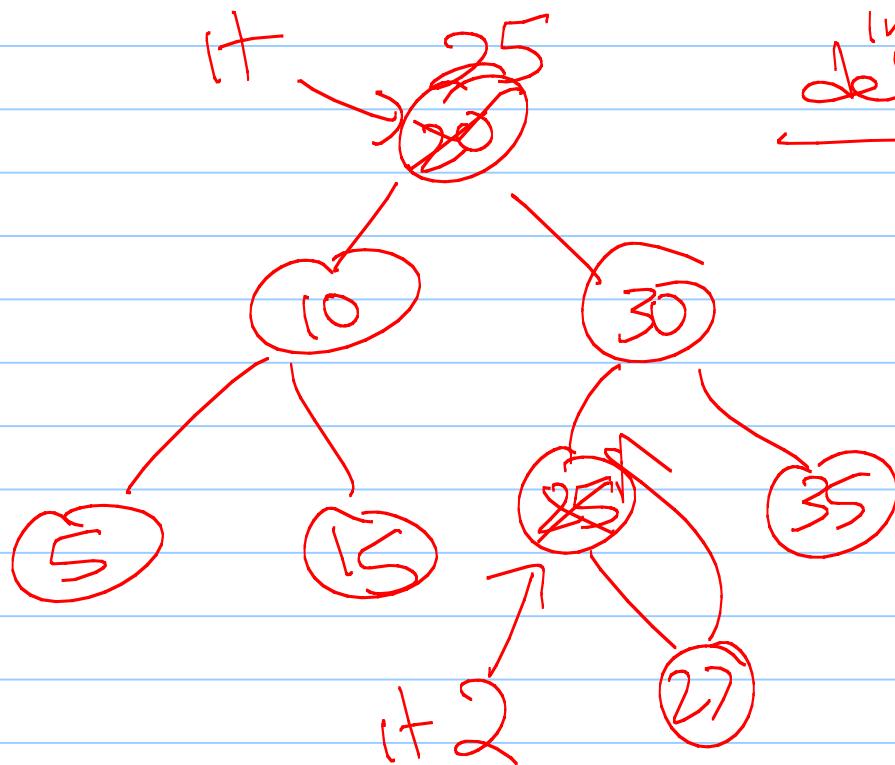
// remove node

remove a leaf - easy!

Case 2: remove(15)



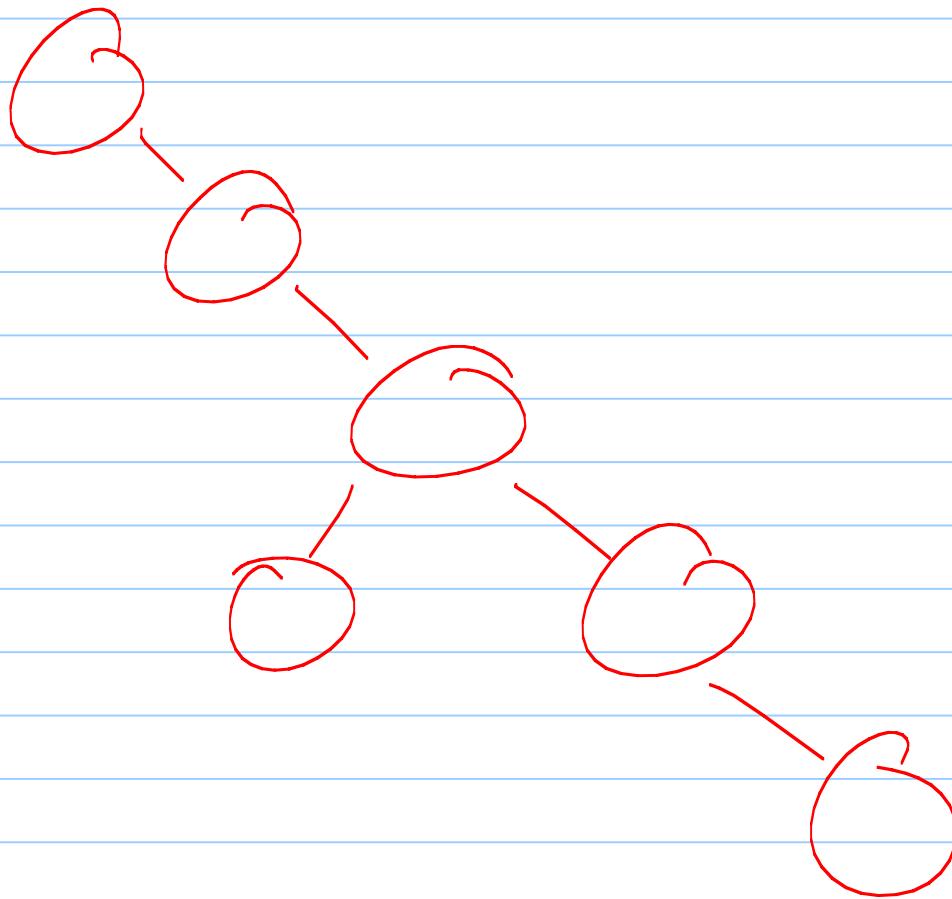
Case 3: 2 children : find next node
in an in-order traversal
(in sorted order)



copy *it 2 into *it
& delete & promote right child
of it 2

only has
(at most) 1
child - on right

$it + 2 = it + + j$

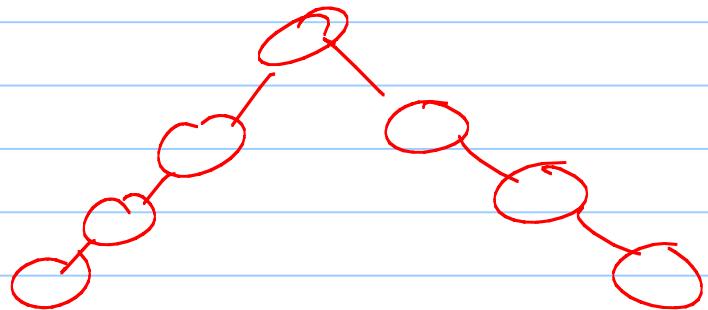


Recap : BST

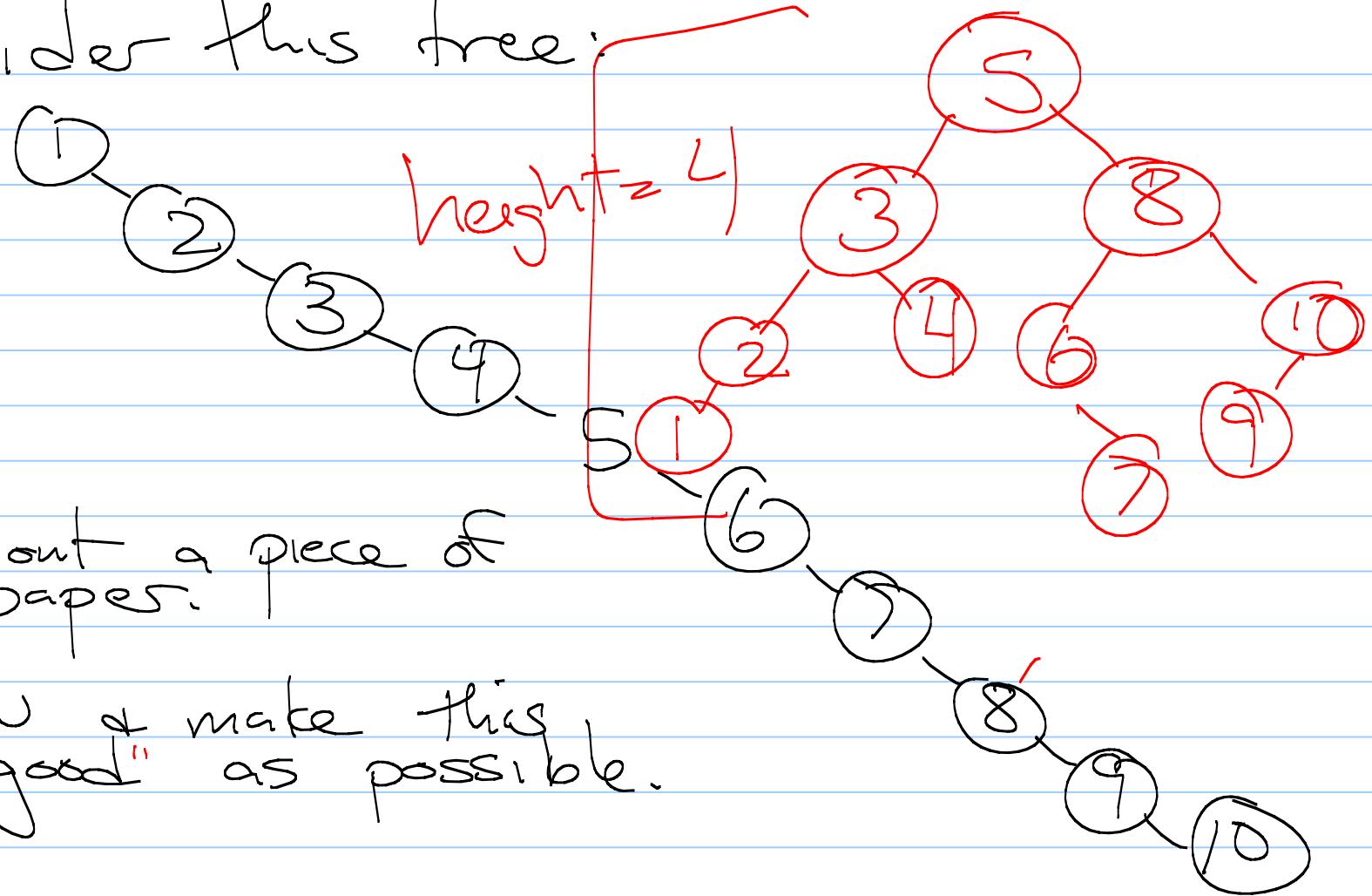
Runtimes:

insert
remove
find } worst case, need
 to travel from
 root to some leaf
 ($\rightarrow O(\text{height})$)

here: $O(n)$



Consider this tree:



Take out a piece of paper.

Redraw & make this as "good" as possible.

What did you do?

Balanced Binary Search Tree

- Red-black trees $\leftarrow 1.2 \log n$

- Splay Trees

- AVL trees

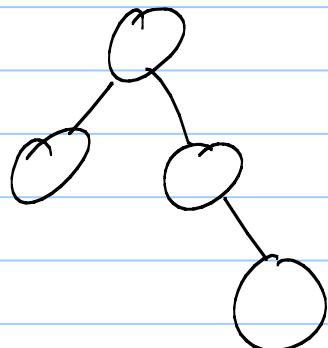
$\leftarrow 1.4 \log n$

Goal of all: $O(\log n)$

AVL Trees

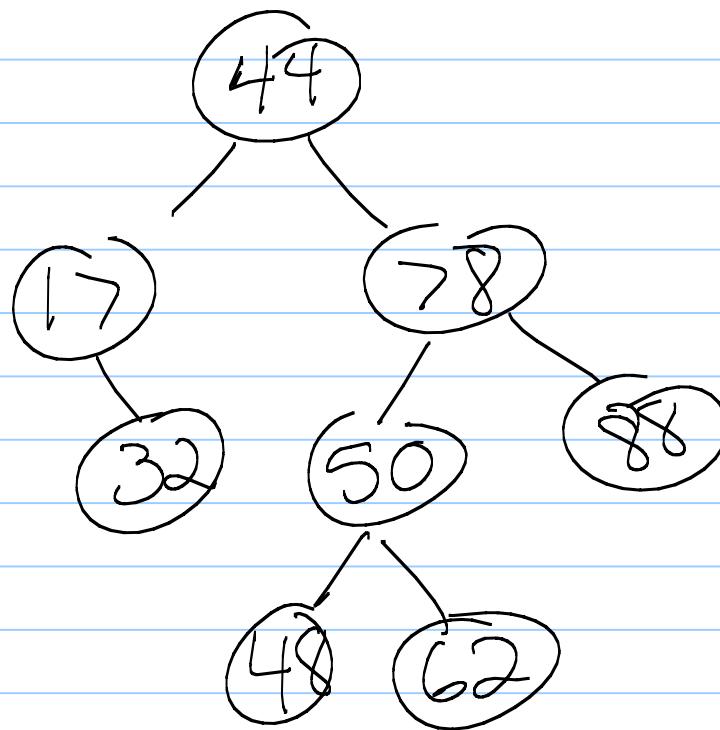
Height - Balance Property :
For every node of T , the heights of the children differ by at most 1.

$$\Rightarrow \max \text{ height} \leq$$



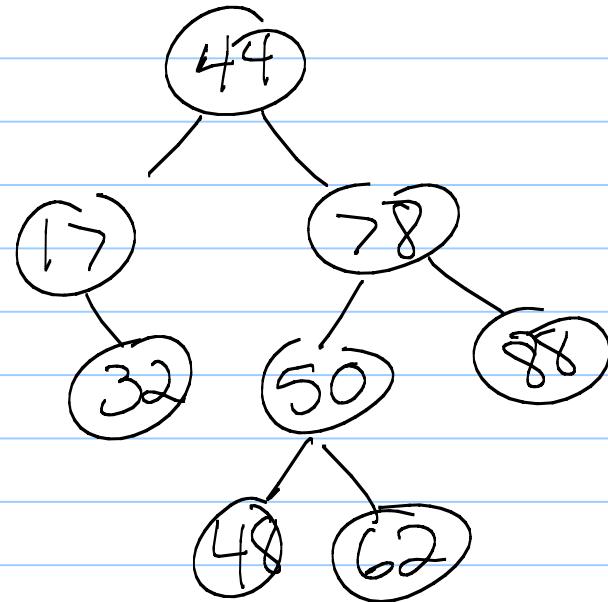
(How do we calculate height again?)

Ex:



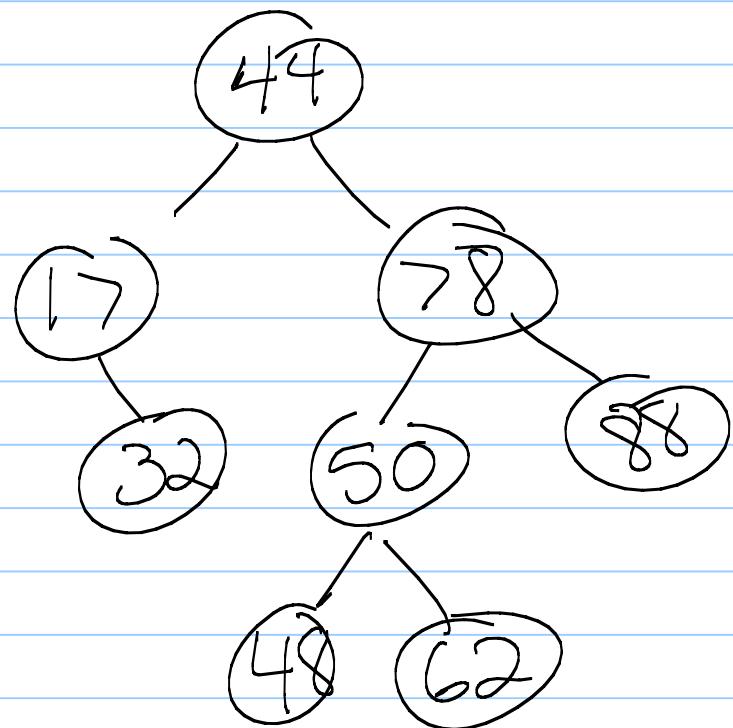
Now: How can we mess this up?

(In other words,
how can the
height change?)



Insert:

insert(54)

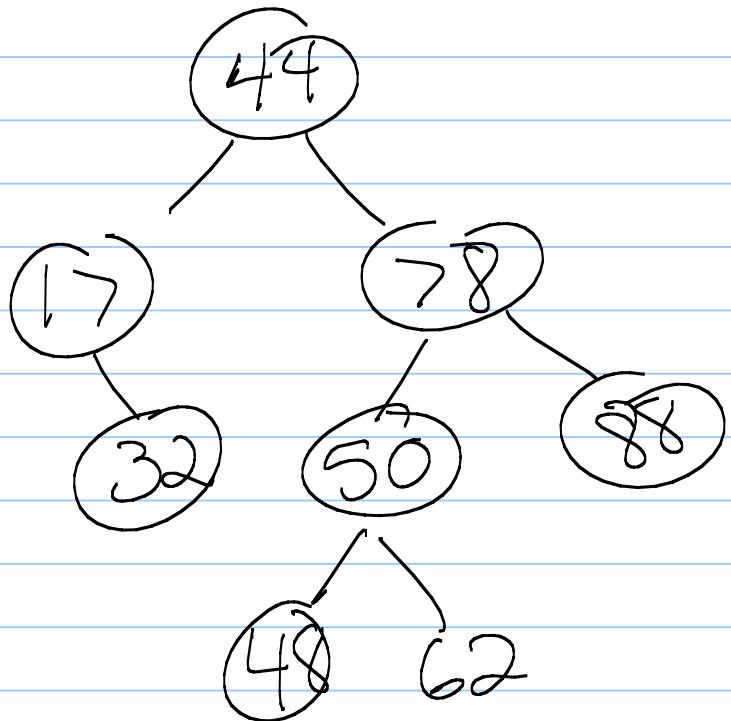


So: consider the lowest node which does not satisfy height-balance property — call this z .

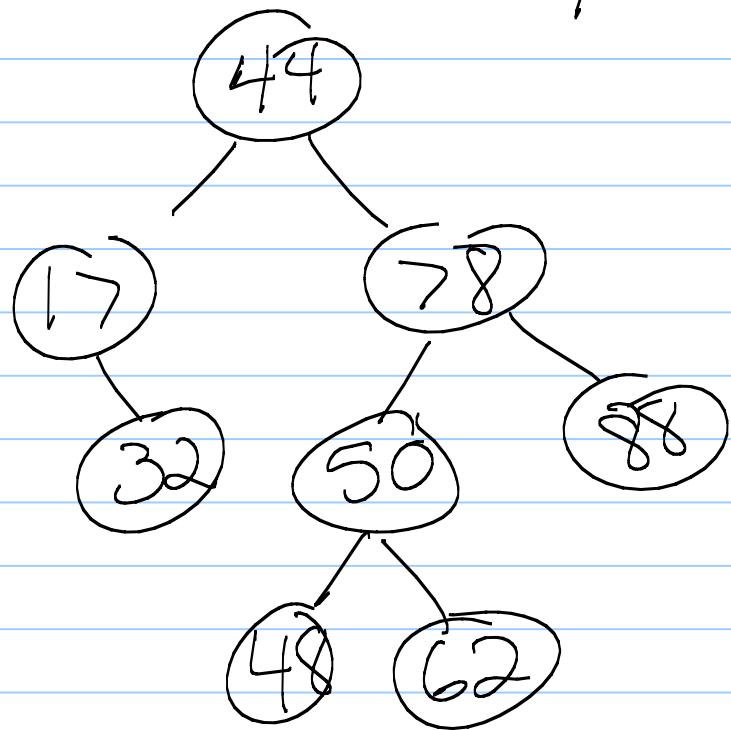
Let y be z 's child with larger height.

Let x be y 's child with larger height.

Now — fix it!



What did you do?



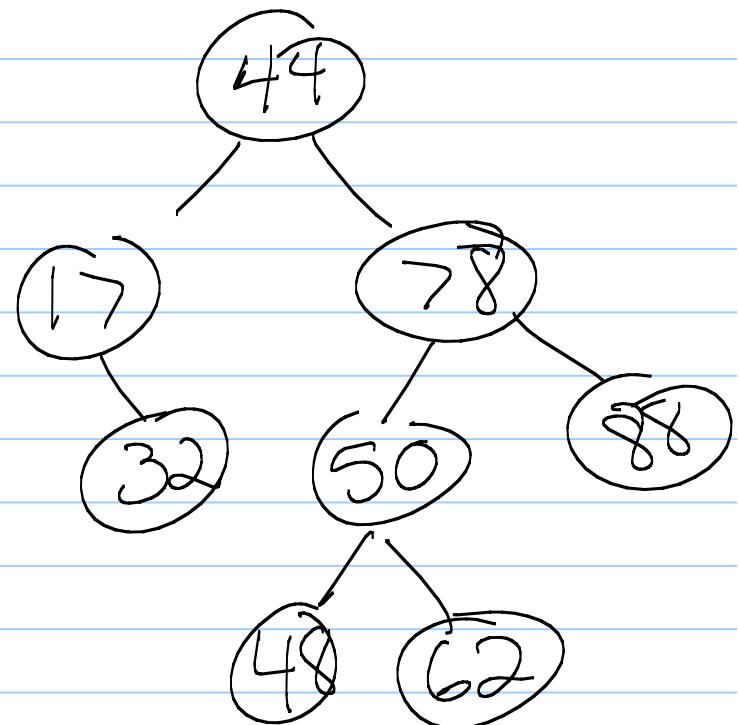
Another - insert 49

So: consider the lowest node which does not satisfy height-balance property - call this z .

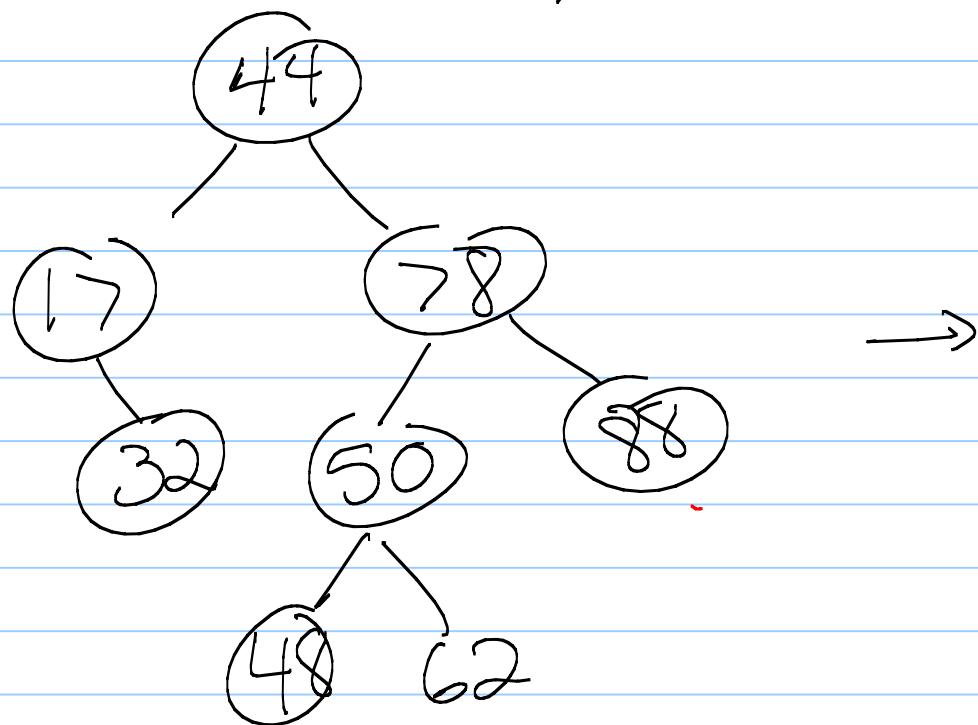
Let x be z 's child with larger height.

Let y be x 's child with larger height.

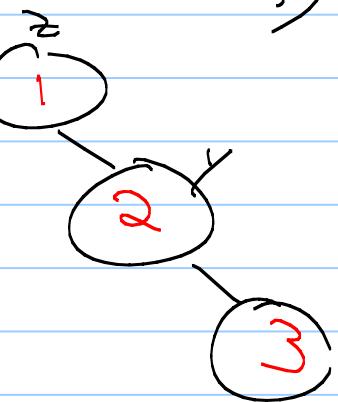
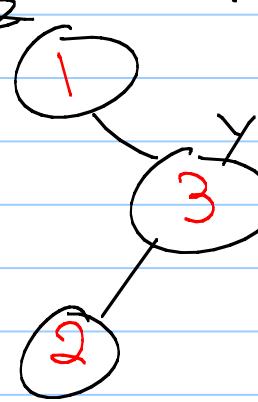
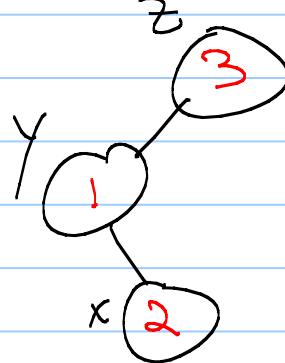
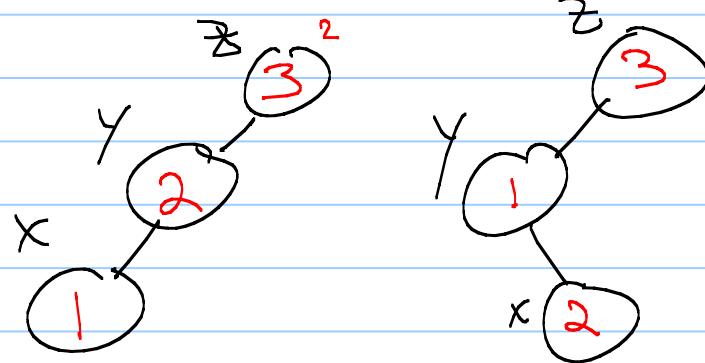
Now - fix it!



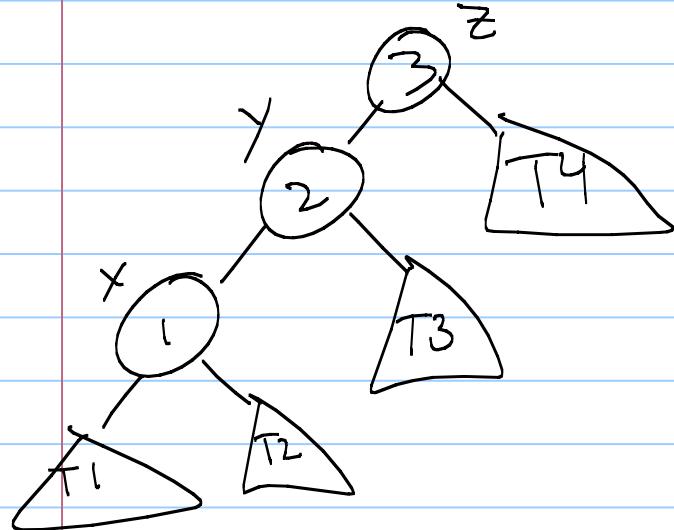
What did you do?



Generalize - Consider x, y , & z . How can we restructure?
(Hint: What is inorder traversal of these in each case?)

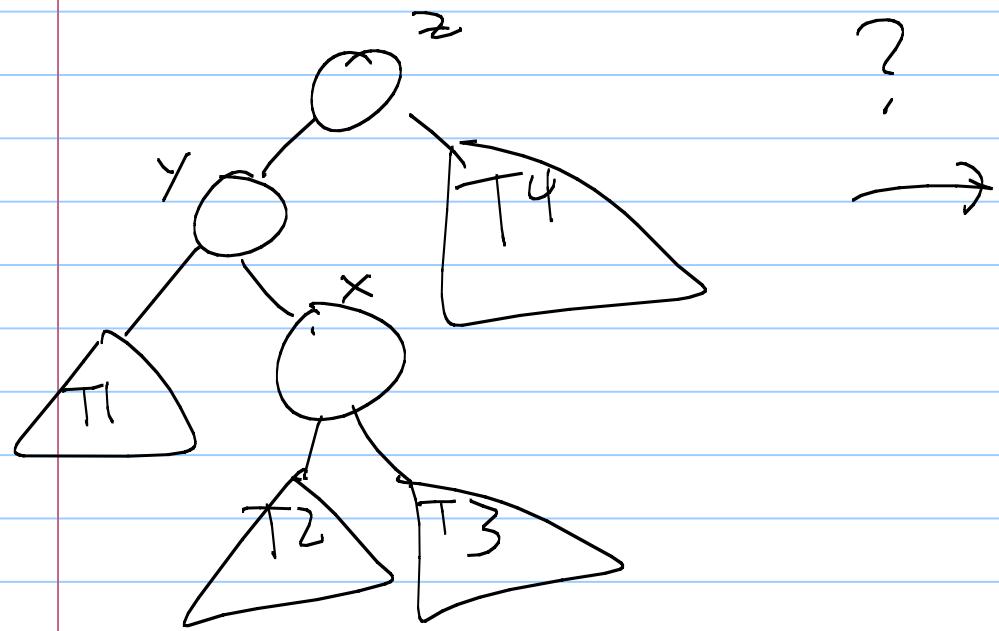


Actual picture:



Where do the subtrees go??

Another



Any way you do this, "2" becomes
the root of the new subtree,
with "1" to the left & "3" to
the right!

What about T1, T2, T3, + T4?

Key operation: Pivot