

CS 314 - Algorithms

Announcements

- Syllabus posted
- HW0 is out - due wed.

Pre reqs

- Discrete Math (or equivalent)
- Data Structures

What is an algorithm?

- Problem Solving
- A set of steps to solve a problem.
- Transform data into information.
- Finite set of steps

What is a program?

- Implementation of an algorithm.
- May combine algorithms.
- Program expresses algorithm.
- Syntax matters!

What is the difference?

Algorithms are much more general.

Here, we care about algorithms.

3 parts to every question:

- ① Algorithm
- ② Runtime analysis
- ③ Proof of correctness

This week: discrete math boot camp!

What did you learn in 135 (or 266)?

- ① Induction !!
- ② Group Theory
- ③ Graph Theory
- ④ Counting
- ⑤ Sorting
- ⑥ Sequences / Summations
- ⑦ Proofs
- ⑧ Logic
- ⑨ Big-O

Runtime:

What is big-O?

Worst case running time

Sorting: $O(n^2)$ (quicksort)
 $O(n \log n)$ (mergesort)

Def:

Let f & g be functions from \mathbb{Z} to \mathbb{R} .
We say $f(x)$ is $O(g(x))$ if

there exist constants $C > 0$ and $N_0 > 0$
such that for all $n > N_0$,

$$f(n) \leq C \cdot g(n)$$

Ex: Show that $x^2 + 2x + 1$ is $O(x^2)$.

proof:

For all $n \geq 1$, $n^2 \geq n^2$ -

For all $n \geq 1$, $2n^2 \geq 2n$ -

For all $n \geq 1$, $n^2 \geq 1$ -

$$\begin{aligned} \text{So } x^2 + 2x + 1 &\leq x^2 + 2x^2 + x^2 \\ &= 4x^2 \end{aligned}$$

Let $C = 4$, then $N_0 = 1$ and

$$x^2 + 2x + 1 \leq C \cdot x^2$$

□

Thm: Let $f(x) = \sum_{i=0}^n a_i x^i = a_n x^n + a_{n-1} x^{n-1} + \dots + a_1 x + a_0$
for constants $a_i \in \mathbb{R}$.
Then $f(x) = O(x^n)$.

(Idea: Let $c = |a_n| + |a_{n-1}| + \dots + |a_1| + |a_0|$)

Other functions

① 2^n versus n^n :

② $\log n$ versus n :