

# Math 135 - Recurrences + Recursion

Note Title

3/3/2010

## Announcements

- HW5 is posted  
due Wed after break
- Scholarships for Math/CS majors  
(for sophomores + juniors)  
applications are in main office (OS RH?)  
(due Wed after break)

# Recurrence Relations

- Use them to model counting problems
- Useful for runtime analysis of recursive algorithms (more next time)
- To solve:
  - unrolling
  - induction
  - more advanced techniques  
such as Master theorem & characteristic eqn method (after break)

Example:

Consider a sequence of numbers:

$a_0$	$a_1$	$a_2$	$a_3$	$a_4$	...
0	2	4	6	8	

Closed form  
 $a_n = 2 \cdot n$

Inductive Dfn:  
 $a_0 = 0$   
 $a_n = a_{n-1} + 2$

Our goal will be  
to figure out  
how to change an  
induction/recursive dfn  
into a closed form.

Another example:

$f_n$ :  $f_0$   $f_1$   $f_2$   $f_3$  - -  
0 1 1 2 3 5 8 13 21 34 ...

Recursive Dfn?

$$\begin{aligned} f_n &= f_{n-1} + f_{n-2} \\ f_0 &= 0 \\ f_1 &= 1 \end{aligned}$$

Closed form?

$$f_n = \frac{(1 + \sqrt{5})^n - (1 - \sqrt{5})^n}{2^n \sqrt{5}}$$

Give recursive defn of:

①  $F(n) = n!$

$\rightarrow F(n) = F(n-1) \cdot n$

②  $A(n) = a^n$

$A(n) = a A(n-1)$

$G(1) = 1$

$G(n) = n \cdot G(n-1)$

unrolling  $= n \cdot (n-1) G(n-2)$   
 $= n(n-1)(n-2) G(n-3)$

⋮

$= n(n-1) \dots 2 \cdot G(1)$

$= n(n-1) \dots 2 \cdot 1$

$= n!$

# Compound interest

- $P_0 = \$1,000$  (initial investment)
- make 6% interest per year

How much will you have after  $n$  years?

Recursive Dfn:  $P_0 = 1,000 \leftarrow$

$$P_n = P_{n-1} + (.06)P_{n-1} = 1.06 P_{n-1}$$

↑  
amount  
from last year

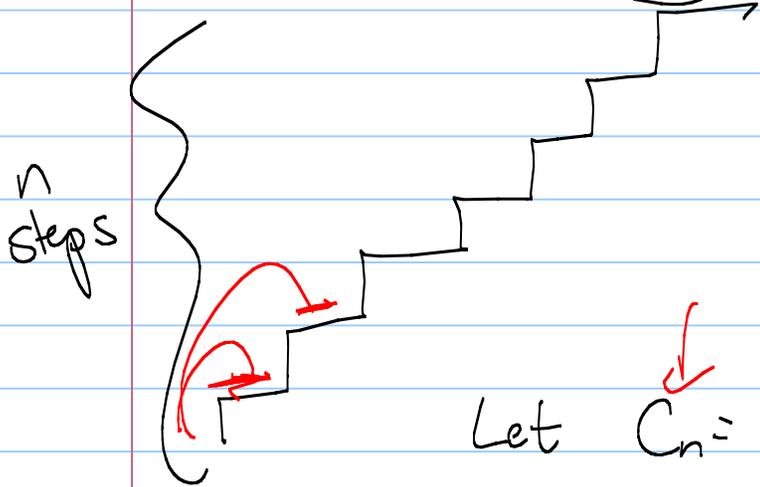
↑  
interest

Unrolling:

Closed form:  $P_n = (1.06) P_{n-1} = (1.06)(1.06) P_{n-2} = (1.06)^2 P_{n-2}$   
 $= (1.06)^2 [1.06 \cdot P_{n-3}] = (1.06)^3 P_{n-3}$

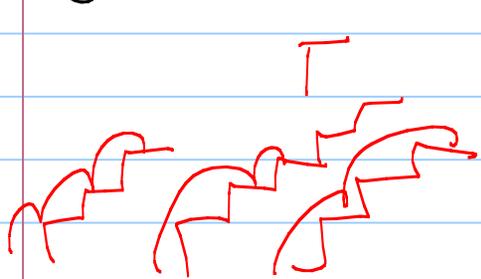


# Stair stepping



If I can take stairs 1 or 2 at a time, how many different ways are there to climb the stairs?

Let  $C_n = \#$  ways to climb  $n$  stairs



$$\begin{aligned} C_1 &= 1 \\ C_2 &= 2 \\ C_3 &= 3 \end{aligned}$$

either I go 2, or 1+then 1

$$C_n = C_{n-1} + C_{n-2}$$

Think recursively!

First step: We could take 1 or 2 stairs.

- Steps we take 1 stair,  
 $C_{n-1}$

- Steps we take 2 stairs,  
 $C_{n-2}$

$$\Rightarrow C_n = C_{n-1} + C_{n-2}$$

Base cases:  $C_1 = 1, C_2 = 2$

Bit strings with no 2 consecutive 0's.

Ex: 1101111  
101010  
~~10011~~

Let  $b_n = \#$  of bit strings w/ no 2 consecutive 0's of length  $n$

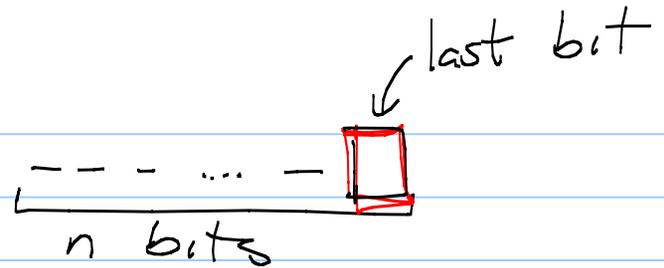
Base cases  $\begin{cases} b_1 = 2 \\ b_2 = 3 \end{cases}$

$\frac{0}{11}$  or  $\frac{1}{10}$   $\frac{0}{01}$   ~~$\frac{0}{00}$~~

Recursive def:  $b_n = ? = b_{n-1} + b_{n-2}$

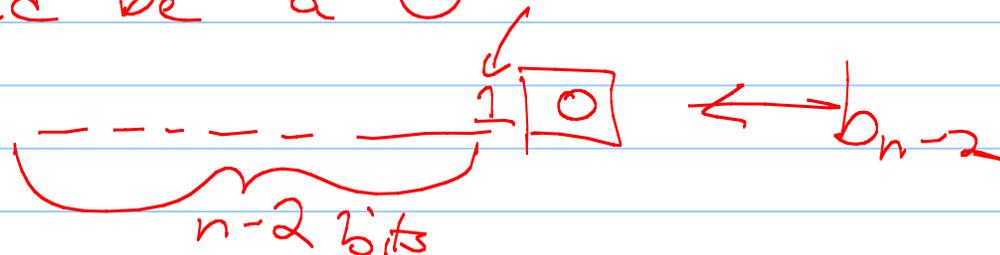
---0---

Consider the last bit:

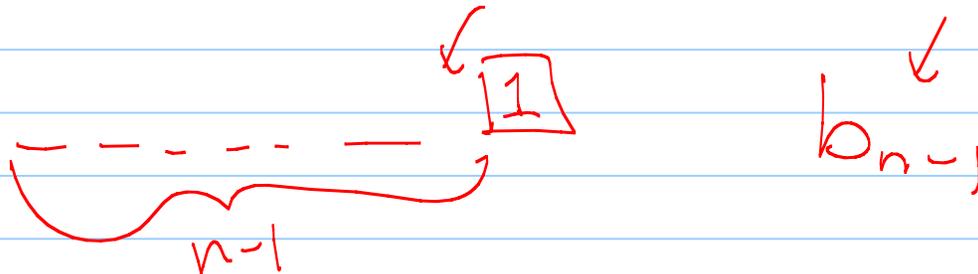


What could it be?

Case 1: could be a 0



Case 2: could be a 1



## Recursively defined sets

Consider an inductive def for a set:

- Base step:  $3 \in S$
- Recursive step: If  $x \in S$  and  $y \in S$ , then  $x+y \in S$ .

So what are some elements of  $S$ ?

$\{3, 6, 9, 15, 12, \dots\}$

Guess? multiples of 3

Claim:  $S = \{ \text{positive integers divisible by } 3 \}$

pf: How do we show 2 sets are equal??

1:  $A \subseteq S$

Induction on elements of  $A = \{3n \mid n \in \mathbb{Z}^+\}$

BC: consider smallest pos. int  $\downarrow$  by 3, = 3

is  $3 \in S$ ?

yes, by base step.

IH:  $\forall k < n, 3k \in S$

IS:  $3n = 3(n-1) + 3$

$\uparrow$   
in  $S$  by IH

$\uparrow$  in  $S$  by B.C.

So by recursive defn,  $3(n-1)+3$  is also in  $S$   $\square$

② Show  $S \subseteq A$

Given any element  $x \in S$ , show  $x \in A$ .

exercise!

(also in text)