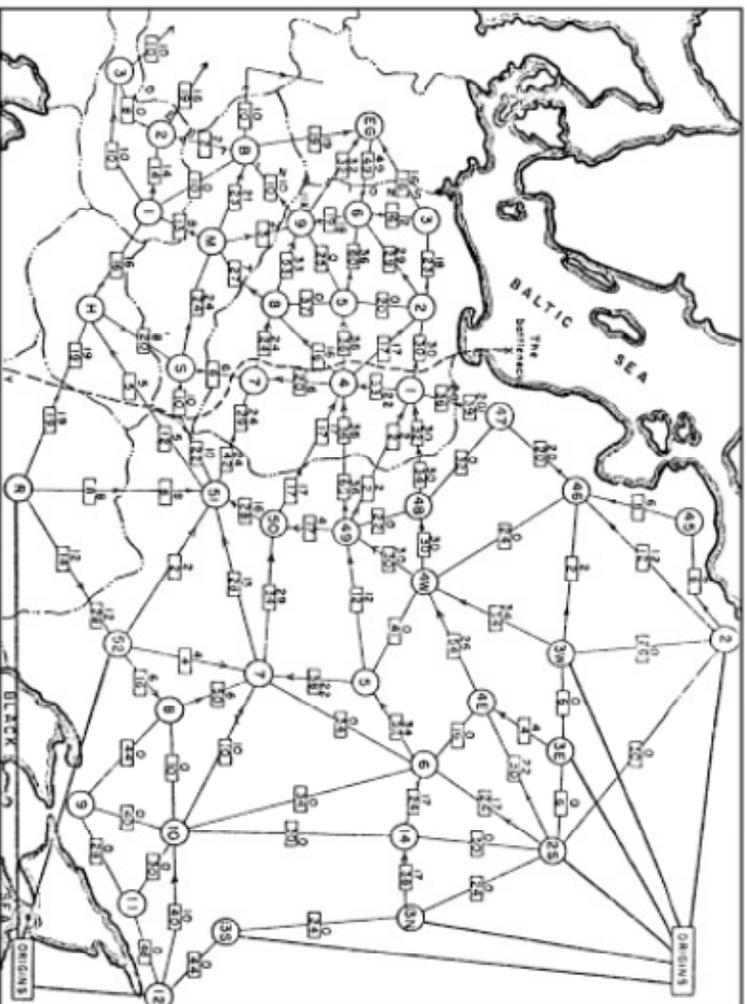


# CS314 - Network Flow

## Announcements

- Midterms graded: average 35.4/60  
(don't worry, will curve!)
- Boering Scholarships
- HW due Friday
- Office hours tomorrow: 11-12

Last week : Network Flow  
(classified report from 1950's)

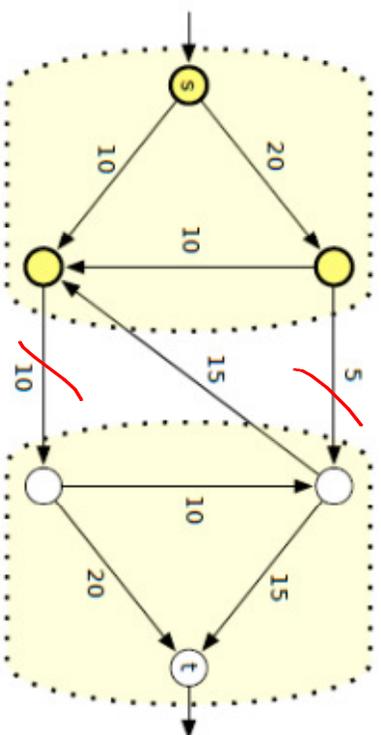
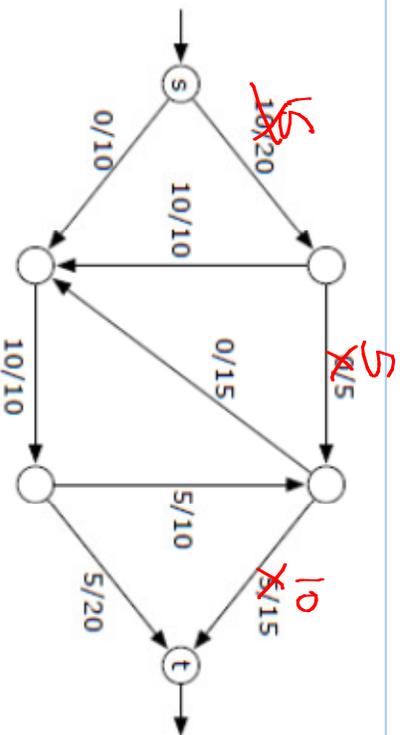


Harris and Ross's map of the Warsaw Pact rail network

Flows

vs.

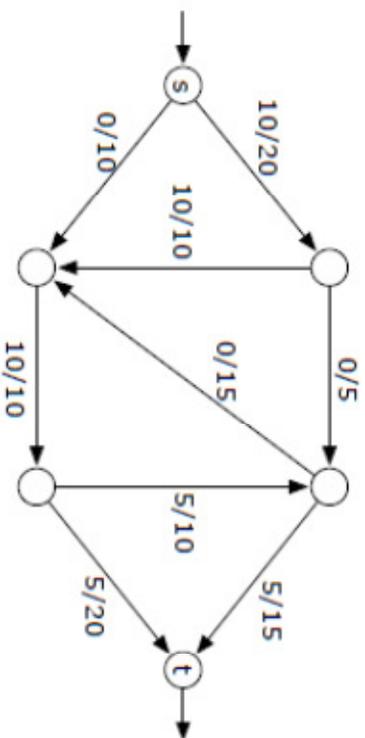
Cuts



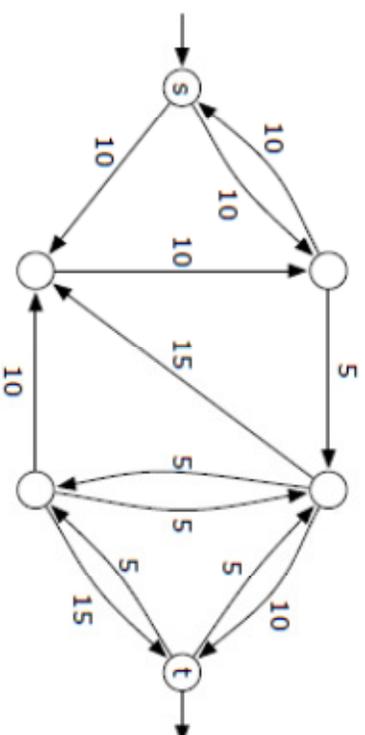
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Key Thm: Max Flow = min cut

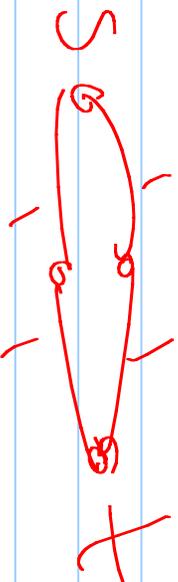
Tool: The residual graph



A flow  $f$  in a weighted graph  $G$  and the corresponding residual graph  $G_f$ .



FF:  $D(m, F)$   
 value of flow



## Today: applications

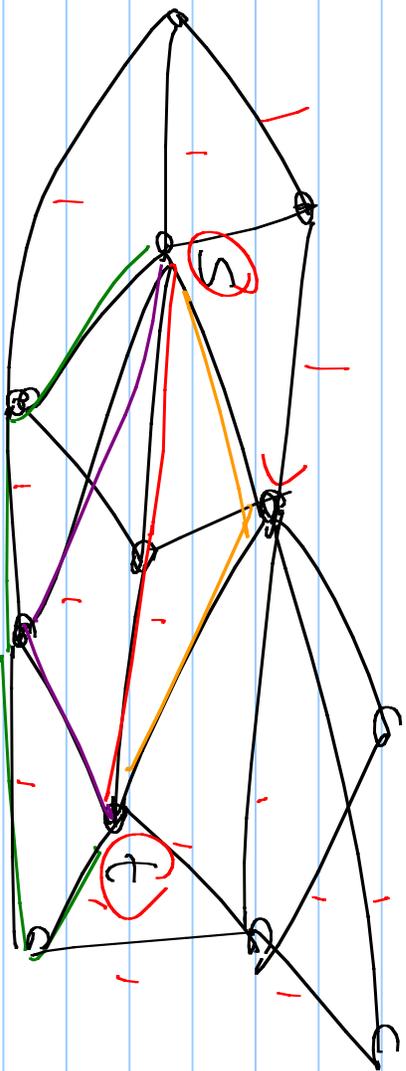
Max flow is useful in many settings.

Given a problem, the main steps are:

- ① Turn it into a "Flow" graph
- ② Plug in F-F  $\rightarrow$  get a runtime
- ③ Correctness  
Flow  $\Leftrightarrow$  solution to problems

## Example: Disjoint Paths

Goal: Find the number of edge disjoint paths in a graph. (between 2 vertices).



(directed)

How?

Put capacity  $\frac{1}{F}$  on every edge

Runtime:

# Correctness:

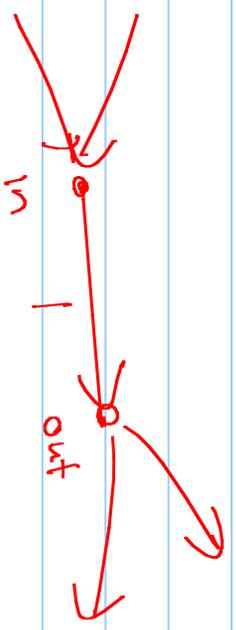
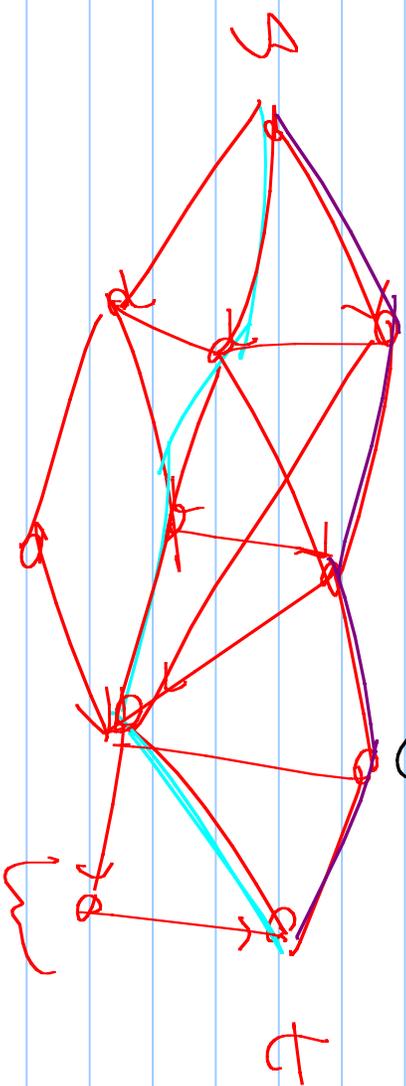
$k$  paths  $\Leftrightarrow$  Flow of size  $k$

$\Rightarrow$  Original  $G$  has  $k$  paths.  
push 1 unit of flow on each path.

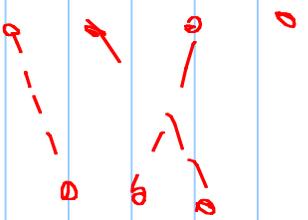
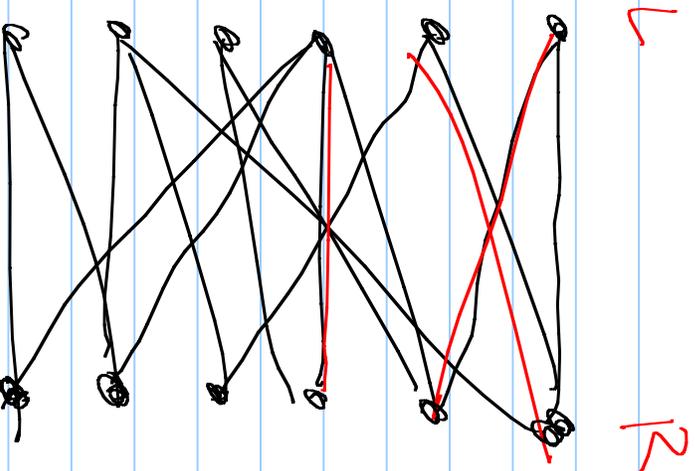


$\Leftarrow$ : If flow of size  $k$  exists, then  $G$  has  $k$  paths.  
Starting at  $s$ , follow edge of length 1. Must be another edge out of flow 1. Continue until hit  $t$ . (repeat  $k$ -times)

What about vertex disjoint paths?



# Example: Bipartite Matching



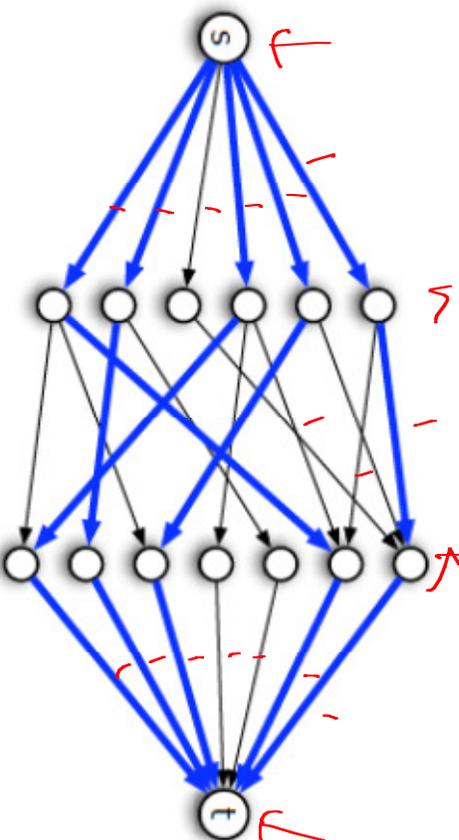
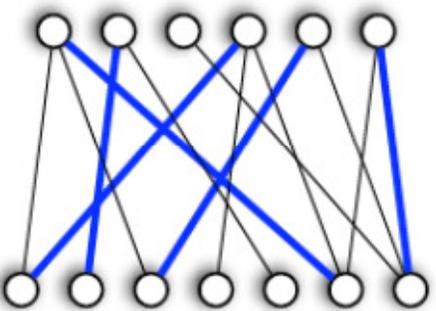
Why?

Scheduling: match people to  
objects

etc...

How?

Flows!



Construct a new graph:  $\# \text{ edges} \leq n/k$

Runtime:

$O(m \cdot F)$  (FF)

||

$O(n \cdot k \cdot \min(n, k))$

# Correctness:

matching  $\Leftrightarrow$  Flow

matchings in original graph: for each  $e=uv$  in matching, send 1 flow along  $s \rightarrow u$ , 1 along  $u \rightarrow v$ , 1 along  $v \rightarrow t$

do this for each edge in matchings  $\Rightarrow$  get a flow



## Another: Assignment Problems

$N$  doctors at a hospital  
 $K$  vacation days

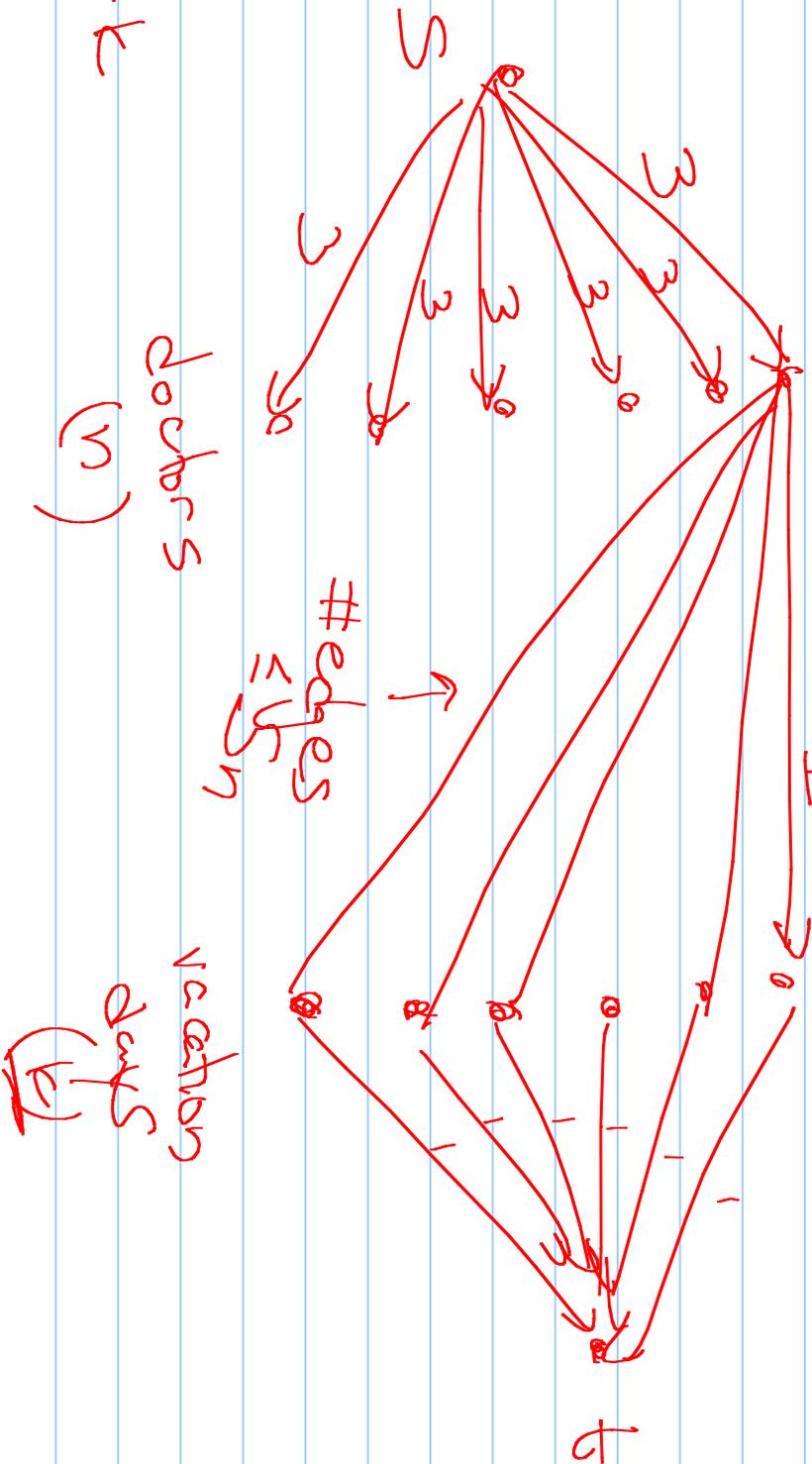
Need:

- a doctor scheduled on every vacation day
- no doctor scheduled on more than 3 vacation days
- each doctor submits a list of  $\geq 5$  vacation days they are available to work.

Q: Is there a feasible schedule?

Q: Is there a flow of size  $k$ ?

Build  $G$ :  
 put edge in  $G$  if  
 doctor not day



Sink

vacation days (k)