

# Math 135 - Algorithms (3.1)

Note Title

2/24/2010

## Announcements

- HW is due on Wed. the 20<sup>th</sup>
- On midterm problem with #10  
Resubmit #10 on Friday

#8 Show  $\sqrt{2}$  is irrational.

Suppose  $\sqrt{2}$  is rational.

$$\Rightarrow \sqrt{2} = \frac{p}{q}, \quad p, q \in \mathbb{Z} \text{ and } q \neq 0.$$

$$\text{then } (\sqrt{2})^2 = \left(\frac{p}{q}\right)^2$$

$$\text{so } \sqrt{2} = \frac{p^2}{q^2}, \quad \frac{p^2}{q^2} \text{ is a rational \#}$$

↳ Since  $\sqrt{2}$  is irrational.

11a

$f(S) = \#$  of 1 bits  
is onto:

Take any  $n \in \mathbb{N}$  and show  
a bitstring  $S$  so that  $f(S) = n$

Let  $S$  be the string with  $n$  1's  
& no zeroes.

$$S = \underbrace{1 \dots 1}_{n \text{ times}}$$

$$A \times B = \{ (a, b) \mid a \in A, b \in B \}$$

#5) Suppose 2 sets with  $A \times B = \emptyset$ .

$$A = \emptyset$$

$$B = \{1, 2\}$$

$$A \times B = \emptyset$$

#9) IH:  $\sum_{k=0}^{n-1} 2^k = 2^n - 1$

IS:  $\sum_{k=0}^n 2^k = \underbrace{\sum_{k=0}^{n-1} 2^k}_{= 2^n - 1} + 2^n = 2^n + 2^n - 1 = 2 \cdot 2^n - 1 = 2^{n+1} - 1$

## Algorithm (Section 3.1)

A set of instructions for solving a problem  
(NOT necessarily a program!)

Examples :

- tying a shoe
- driving
- walking
- recipes

We often use pseudocode to write down computer algorithms.

Common programming concepts:

- if statements
- loops
- variables
- functions or procedures
- input/output

Ex: Pseudocode to find the maximum element  
in a sequence  $a_1 \dots a_n$

**FINDMAX**( $a_1, a_2, \dots, a_n$ ):

```
max :=  $a_1$ 
for  $i := 2$  to  $n$ 
  if  $max < a_i$ 
     $max := a_i$ 
return max
```

Two important  
structures:

for loops  
& if statements

# Searching

Suppose I give you a list of numbers  $a_1, \dots, a_n$   
and ask if  $x \in \{a_1, \dots, a_n\}$ .  
How would you check?

for loop to check each element



while (bool)

→ if bool is true,  
execute the instructions  
inside the loop

LINEAR SEARCH( $x, a_1, \dots, a_n$ ):

```
i := 1
while (i ≤ n and x ≠ ai)
  i := i + 1
if i ≤ n
  location := i
else
  location := 0
return location
```

~~i = x~~, n = 5  
3

10 12 -1 3 5

x = -1

i = 1 ?

Another Search Strategy:

Ex: Take out your book & open it to page 171.

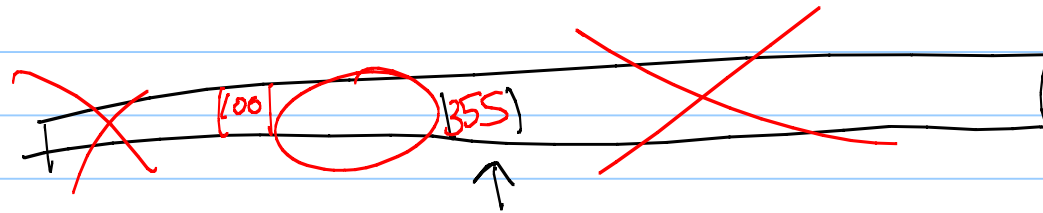
How does your algorithm to do this differ from the linear search algorithm?

binary search - check the middle

(assume the list is sorted!)

When searching in a sorted list, we can do a faster search called binary search.

- Compare to middle element of list.



- If that element is bigger than  $x$ ,  $x=17$   
search in the left half

- If that element is smaller than  $x$ ,  
search in the right  
(pseudocode in book)

# Sorting:

Fundamental CS problem:

Given a list of  $n$  things, put them in order

How?

Comparing elements, & swap  
if out of order

## Bubble Sort:

Compare adjacent elements + switch them  
if in wrong order

3 2 4 1 5

2 3 4 1 5

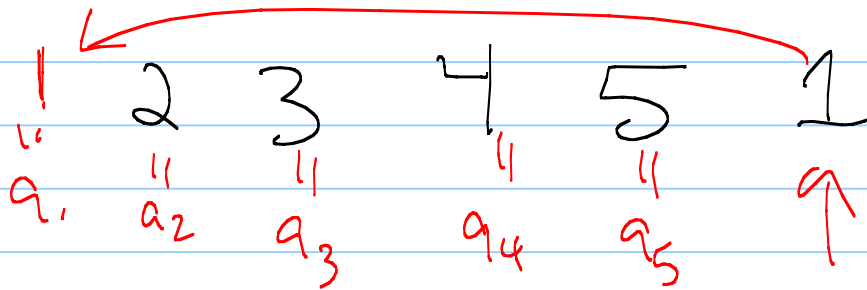
2 3 1 4 5

## Pseudocode

```
BUBBLE SORT ( $a_1 \dots a_n$ ):  
  for  $i := 1$  to  $n-1$   
    for  $j := 1$  to  $n-i$   
      if  $a_j > a_{j+1}$   
        swap  $a_j$  and  $a_{j+1}$ 
```

# Insertion Sort

- If first  $i$  items are sorted, take  $(i+1)^{th}$  and put it in correct spot.



## Pseudo code

INSERTION SORT ( $a_1 \dots a_n$ ):

for  $j := 2$  to  $n$

$i := 1$

while  $a_j > a_i$

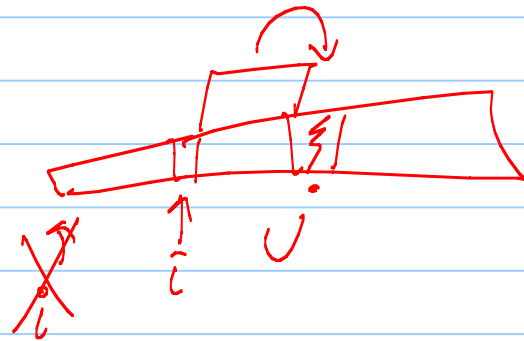
$i := i + 1$

temp :=  $a_j$

for  $k := j$  to  $i + 1$

$a_k := a_{k-1}$

$a_i := \text{temp}$





# Complexity of Algorithms

Comparing which algorithms are "better" can be tricky. U

Issues:

- time them on a machine
- is input good?
- different on different machines

So:

We define complexity in terms of the number of operations.

Usually, an operation is:

- add 2 things (or subtract or multiply)
- compare 2 things
- set a variable equal to something

atomic operations

But still - how do we compare?

We just saw 2 searching algorithms,  
linear search & binary search.

One is not always better.

Find(36):

36 40 58 100 101 125

Linear search: 1  
Binary search: 3

Find(36):

1 11 25 36 41 42 100

Linear search: 4  
Binary search: 1

So how can we compare worst case performance?

~~✱~~ worst case performance

Ex: What is worst case complexity of FindMax?

$$1 + (n-1)(2) + 1$$
$$= 2n - 2 + 2 = 2n$$

$n-1$

1

```
FINDMAX( $a_1, a_2, \dots, a_n$ ):  
  max :=  $a_1$   
  for  $i := 2$  to  $n$   
    if max <  $a_i$   
      max :=  $a_i$   
  return max
```