

Math 35 - Undecidable Problems

Note Title

3/1/2010

Announcements

- Midterm grades will be posted ASAP
- Next HW will be posted later today
(due next Wed. or Fri.)
- Second exam will be in ~3 weeks

Last time: Algorithm Complexity

We use big- O . (x count primitive operations)

Why?

- processor independent
- worst case
- language independent

We saw $O(1)$ to $O(2^n)$
 $O(n)$, $O(n^2)$

The Halting Problem

Undecidable problems

(end of 3.1)

Q: Can we write a program which accepts as input another program & input, then decides if the program will run forever or halt on that input.

(So if it contains infinite loop, will run forever, for example, & our program will say that.)

Note: Our program can't just run
the input \cup program.

Why?

If I simulate the other program
& it runs forever, then
so do I.

Thm: The halting problem is undecidable.
(that is, no program to solve it,
can exist!)

pf: by contradiction

Assume $H(P, I)$ takes P (a program)
& I (an input) & outputs "halts"
or "runs forever".

Any program is just a string of bits.

So H is also a program
 $\rightarrow H(H, I) \cup H(P, P)$

(continued) Design a program that uses H ,
called K .
 $K(P)$ calls $H(P, P)$.
if $H(P, P)$ says "loop forever", K will
halt.
If $H(P, P)$ says "halt", K will loop
forever.

What does $K(K)$ do?
depends on $H(K, K)$.
if it says "loop forever", $K(K)$
will halt.
if it says "halt", then $K(K)$
loops forever.
either way, reached a contradiction. @

Recurrence Relations (4.3-4.4)

- Use them to model counting problems
- Useful for runtime analysis of recursive algorithms (more next time)
- To solve:
 - unrolling
 - • induction
 - more advanced techniques
such as Master theorem & characteristic eqn method (after break)

Example:

Consider a sequence of numbers:

a_0	a_1	a_2	a_3	a_4	...
0	2	4	6	8	

Closed form

$$a_n = 2 \cdot n$$

Recursive

~~Inductive~~ Dfn:

$$a_0 = 0$$

$$a_n = a_{n-1} + 2$$

Another example: Fibonacci numbers

	f_0	f_1	f_2	f_3	f_4	f_5	f_6	f_7	f_8	f_9	...
f_n :	0	1	1	2	3	5	8	13	21	34	...

$$f_0 = 0$$
$$f_1 = 1$$

Recursive Defn?

$$f_n = f_{n-1} + f_{n-2}$$

Closed form?

$$f_n = \frac{\left(\frac{1+\sqrt{5}}{2}\right)^n - \left(\frac{1-\sqrt{5}}{2}\right)^n}{\sqrt{5}}$$

Give recursive defn of:

$$\textcircled{1} F(n) = n! \quad F(1) = 1$$

$$F(n) = n \cdot F(n-1)$$

$$n! = n (n-1)!$$

$$\textcircled{2} A(n) = a^n$$

$$\downarrow A(n) = a \cdot A(n-1)$$

$$A(n-1) = a^{n-1}$$

Compound interest

- $P_0 = \$1,000$ (initial investment)
- make 6% interest per year

How much will you have after n years?

Recursive Dfn: $P_0 = 1,000$ ←

$$P_n = (1.06) P_{n-1}$$

(Note: In the original image, the entire recursive formula is circled in red. An arrow points from the 1.06 to the P_{n-1} term, and another arrow points from the P_{n-1} term to the P_n term.)

$$P_n = (1.06)^n (1,000) \leftarrow$$

Claim: $P_n = (1.06)^n P_0 = (1.06)^n (1,000)$

pf: (induction on n)

base case: $n=0$ $P_0 = 1000$
 $(1.06)^0 \cdot 1000 = 1 \cdot 1000 = 1000$ ✓

ind hyp: Assume $P_{n-1} = (1.06)^{n-1} \cdot (1000)$

ind step: Consider P_n .

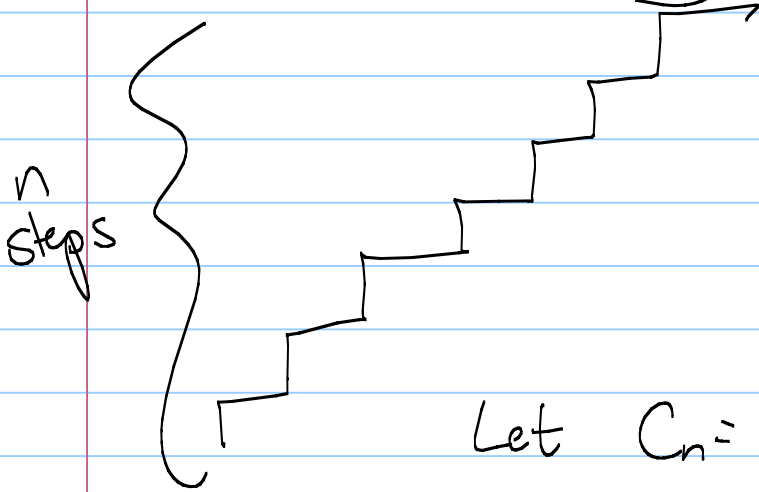
Know $P_n = (1.06) P_{n-1}$ by recursive def.

By IH $= (1.06) (1.06)^{n-1} (1000)$

$= (1.06)^n (1000)$

✓ \square

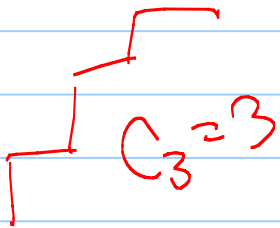
Stair stepping



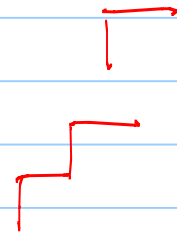
If I can take stairs 1 or 2 at a time, how many different ways are there to climb the stairs?

Let $C_n = \#$ ways to climb n stairs

Base cases:



1+1



or 2

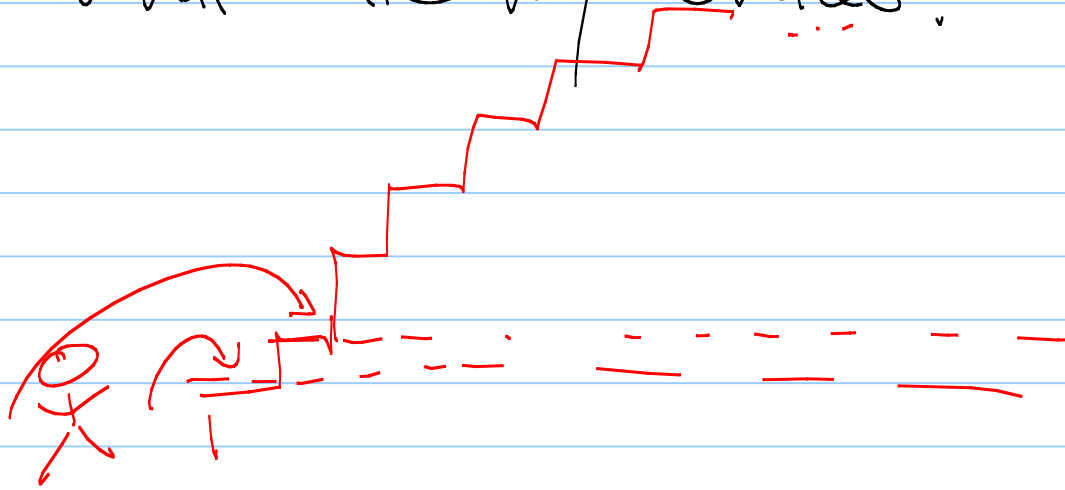
$C_2 = 1$

$C_2 = 2$

Think recursively!

First step: What are my choices?

C_n



$$C_n = C_{n-1} + C_{n-2} = (C_{n-2} + C_{n-3}) + (C_{n-3} + C_{n-4})$$

Fibonacci numbers!